

## Chapter 3

### Exercise 3.1

$v$  has access to no world. So any sentence  $A$  is true at *all* (zero) worlds accessible from  $v$ .

If this seems strange, remember that  $\Box A$  is equivalent to  $\neg \Diamond \neg A$ . And  $\Diamond \neg A$  means that there's an accessible world where  $\neg A$  is true. If there are no accessible worlds, then this is false. So  $\neg \Diamond \neg A$  is true.

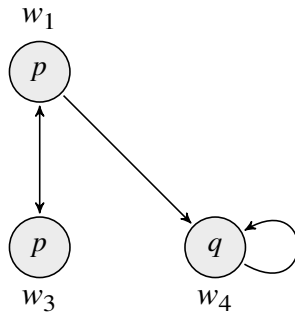
### Exercise 3.2

(a)  $w_1, w_2$ , and  $w_3$ ; (b)  $w_3$ ; (c)  $\neg$ ; (d)  $w_1, w_2$  and  $w_4$ ; (e) all.

### Exercise 3.3

There are infinitely many correct answers for each world. For example:  $w_1 : \Diamond \Box p$ ,  $w_2 : \neg p \wedge \neg q$ ,  $w_3 : \Box p$ ,  $w_4 : \Box q$ .

### Exercise 3.4



### Exercise 3.5

- (a) For example:  $W = \{w, v\}$ ,  $R = \{(w, v), (v, w)\}$ ,  $V(p) = \{v\}$ .  $\Box p \rightarrow \Box \Box p$  is false at  $w$ . (' $R = \{(w, v), (v, w)\}$ ' means that  $R$  relates  $w$  to  $v$  and  $v$  to  $w$  and nothing else to anything else.)
- (b) For example:  $W = \{w, v\}$ ,  $R = \{(w, w), (w, v)\}$ ,  $V(p) = \{w\}$ .  $\Diamond p \rightarrow \Box \Diamond p$  is false at  $w$ .

### Exercise 3.6

For example:  $\Box(p \vee \neg p) \rightarrow (p \vee \neg p)$ .

### Exercise 3.7

By clause (g) of definition 3.2,  $\Box(p \vee \neg p)$  is false at a world  $w$  in a Kripke model only if  $p \vee \neg p$  is false at some world accessible from  $w$ . By clause (d) of definition 3.2,  $p \vee \neg p$  is false at a world only if both  $p$  and  $\neg p$  are false at the world, which by clause (a) means that  $p$  is both true and false at the world. This is impossible. So  $\Box(p \vee \neg p)$  is not false at any world in any Kripke model.

### Exercise 3.8

By definition 3.2,  $\Box p \rightarrow \Diamond p$  is false at a world  $w$  in a Kripke model only if  $\Box p$  is true at  $w$  and  $\Diamond p$  is false at  $w$ . But if  $w$  has access to itself then the truth of  $\Box p$  at  $w$  implies that  $p$  is true at  $w$ , and then  $\Diamond p$  is false at  $w$ . So  $\Box p \rightarrow \Diamond p$  can't be false at any world in any Kripke model in which each world has access to itself.

### Exercise 3.9

Reflexive yes, serial yes, transitive yes, euclidean no, symmetric no, universal no.

### Exercise 3.10

- (a) Suppose  $R$  is symmetric and transitive, and that  $xRy$  and  $xRz$ . By symmetry,  $yRx$ . By transitivity,  $yRz$ .
- (b) Suppose  $R$  is symmetric and euclidean, and that  $xRy$  and  $yRz$ . By symmetry,  $yRx$ . By euclidity,  $xRz$ .
- (c) Suppose  $R$  is reflexive and euclidean, and that  $xRy$ . By reflexivity,  $xRx$ . By euclidity,  $yRx$ .

### Exercise 3.11

It's true that if  $R$  is symmetric and transitive then  $wRv$  implies  $vRw$  which implies  $wRw$ . But this only shows that every world  $w$  that can see some world  $v$  can see

itself. Symmetry, transitivity, *and seriality* together imply reflexivity. Symmetry and transitivity alone do not.

**Exercise 3.12**

- (a) Every world has access only to itself.
- (b) No world has access to any world.

**Exercise 3.13**

You can enter the sentences at [umsu.de/trees](https://umsu.de/trees). To check for K-validity, leave all the checkboxes (for ‘universal’ etc.) empty.

**Exercise 3.14**

You can enter the sentences at [umsu.de/trees](https://umsu.de/trees). To test for K4-validity, check the ‘transitive’ box. To test for D-validity, check ‘serial’. To test for B-validity, check ‘symmetric’. To test for T-validity, check ‘reflexive’.